



# **Background/Review on Numbers and Computers (lecture)**

## **ICS312 Machine-Level and Systems Programming**

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# Numbers and Computers

- Throughout this course we will
  - use **binary** and **hexadecimal** representations of numbers
  - need to be aware of the ways in which the computer stores numbers
- So let us go through a simple review before we start learning how to write assembly code
  - Numbers in different bases
  - Number representation in computers and basic arithmetic
    - More to come later on arithmetic

# Numbers and bases

- We are used to thinking of numbers as written in decimal, that is, in base 10

$$25 = 2 \cdot 10^1 + 5 \cdot 10^0$$

$$136 = 1 \cdot 10^2 + 3 \cdot 10^1 + 6 \cdot 10^0$$

- Each number is decomposed into a sum of terms
- Each term is the product of two factors
  - A digit (from 0 to 9)
  - The base (10) raised to a power corresponding to the digit's position in the number

$$\begin{aligned} 136 &= \dots + 0 \cdot 10^4 + 0 \cdot 10^3 + 1 \cdot 10^2 + 3 \cdot 10^1 + 6 \cdot 10^0 \\ &= \dots 00000136 \end{aligned}$$

- We typically don't write (an infinite number of) leading 0's

# Numbers and Bases

- Any number can be written in base  $b$ , using  $b$  digits
  - If  $b = 10$  we have “**decimal**” with 10 digits [0-9]
  - If  $b = 2$  we have “**binary**” with 2 digits [0,1], which are also called **bits**
  - If  $b = 8$  we have “**octal**” with 8 digits [0-7]
  - If  $b = 16$  we have “**hexadecimal**” with 16 digits [0-9,A,B,C,D,E,F]
- Computers use binary internally
  - It's easy to associate two states to a current
    - Low voltage = 0, high voltage = 1
    - Associating 16 states to a current is more complicated and error-prone
- However, binary is cumbersome
  - The lower the base the longer the numbers!
  - It's really difficult for a human to remember binary
- Therefore we, as humans, like to use higher bases
- **Bases that are powers of 2 make for easy translation to binary**, and thus are particularly useful, and in particular **hexadecimal**

# Binary Numbers

- Counting in binary:

$0_2$                        $0_{10}$

$1_2$                          $1_{10}$

$10_2$                        $2_{10}$

$11_2$                        $3_{10}$

$100_2$                      $4_{10}$

$101_2$                      $5_{10}$

$110_2$                      $6_{10}$

$111_2$                      $7_{10}$

$1000_2$                    $8_{10}$

...

- A binary number with  $d$  bits corresponds to integer values between 0 and  $2^d - 1$

$$\sum_{k=0}^{d-1} 2^k = 2^d - 1$$

- Example:
  - An integer stored in 8 bits has values between 0 and 255
  - $128 + 64 + 32 + 16 + 8 + 4 + 2 + 1 = 255$

# Converting from Binary to Decimal

- We denote by  $XXXX_2$  a binary representation of a number and by  $XXXX_{10}$  a decimal representation
- Converting from binary to decimal is straightforward:

$$\begin{aligned}10010110_2 &= 1 \cdot 2^7 + 1 \cdot 2^4 + 1 \cdot 2^2 + 1 \cdot 2^1 \\ &= 1 \cdot 128 + 1 \cdot 16 + 1 \cdot 4 + 1 \cdot 2 \\ &= 150_{10}\end{aligned}$$

- The rightmost bit of a binary number is called the **least significant bit**
- The leftmost non-zero bit of a binary number is called the **most significant bit**
- If the least significant bit is 0, then the number is even, otherwise it's odd

# Converting from Decimal to Binary

- The conversion proceeds by a series of **integer divisions** by 2, and by recording the remainder of the division
  - Integer division  $a/b$ :  $a = b * q + \text{remainder}$ , where all are integers
- Example: converting  $37_{10}$  into binary
  - Divide 37 by 2:  $37 = 2 * 18 + 1$
  - Divide 18 by 2:  $18 = 2 * 9 + 0$
  - Divide 9 by 2:  $9 = 2 * 4 + 1$
  - Divide 4 by 2:  $4 = 2 * 2 + 0$
  - Divide 2 by 2:  $2 = 2 * 1 + 0$
  - Divide 1 by 2:  $1 = 2 * 0 + 1$
  - Result:  $100101_2$
- **The least significant bit is computed first**
- **The most significant bit is computed last**
- Note that if we continue dividing, we get extraneous leading 0s
  - $\dots 00000100101_2$

# Binary Arithmetic

- Adding a 0 to the right of a binary number multiplies it by 2
  - $10101_2 = 16_{10} + 4_{10} + 1_{10} = 21_{10}$
  - $101010_2 = 32_{10} + 8_{10} + 2_{10} = 42_{10}$
- Adding two binary numbers is just like adding decimal numbers: using a carry

With no previous carry				With a previous carry			
0	0	1	1	0	0	1	1
+ 0	+ 1	+ 0	+ 1	+ 0	+ 1	+ 0	+ 1
= 0	= 1	= 1	= 0	= 1	= 0	= 0	= 1
			C		C	C	C



# Binary Addition

$$\begin{array}{r} \text{c c c c} \\ 1001 \\ + 1111 \\ = 11000 \end{array}$$

$$\begin{array}{r} 9_{10} \\ + 15_{10} \\ = 24_{10} \end{array}$$

$$\begin{array}{r} \text{c} \qquad \qquad \text{c c} \\ 10100110 \\ + 11000011 \\ = 101101001 \end{array}$$

$$\begin{array}{r} 166_{10} \\ + 195_{10} \\ = 361_{10} \end{array}$$

# Counting in Hexadecimal

$$0_{16} = 0_{10}$$

$$1_{16} = 1_{10}$$

$$2_{16} = 2_{10}$$

$$3_{16} = 3_{10}$$

$$4_{16} = 4_{10}$$

$$5_{16} = 5_{10}$$

$$6_{16} = 6_{10}$$

$$7_{16} = 7_{10}$$

$$8_{16} = 8_{10}$$

$$9_{16} = 9_{10}$$

$$A_{16} = 10_{10}$$

$$B_{16} = 11_{10}$$

$$C_{16} = 12_{10}$$

$$D_{16} = 13_{10}$$

$$E_{16} = 14_{10}$$

$$F_{16} = 15_{10}$$

$$10_{16} = 16_{10}$$

$$11_{16} = 17_{10}$$

$$12_{16} = 18_{10}$$

$$13_{16} = 19_{10}$$

$$14_{16} = 20_{10}$$

$$15_{16} = 21_{10}$$

$$16_{16} = 22_{10}$$

$$17_{16} = 23_{10}$$

$$18_{16} = 24_{10}$$

$$19_{16} = 25_{10}$$

$$1A_{16} = 26_{10}$$

$$1B_{16} = 27_{10}$$

$$1C_{16} = 28_{10}$$

$$1D_{16} = 29_{10}$$

$$1E_{16} = 30_{10}$$

$$1F_{16} = 31_{10}$$

$$20_{16} = 32_{10}$$

$$21_{16} = 33_{10}$$

$$22_{16} = 34_{10}$$

$$23_{16} = 35_{10}$$

$$24_{16} = 36_{10}$$

$$25_{16} = 37_{10}$$

$$26_{16} = 38_{10}$$

$$27_{16} = 39_{10}$$

# Converting from hex to decimal

- This is again straightforward

$$\begin{aligned} A203DE_{16} &= 10 \cdot 16^5 + \\ &\quad 2 \cdot 16^4 + \\ &\quad 3 \cdot 16^2 + \\ &\quad 13 \cdot 16^1 + \\ &\quad 14 \cdot 16^0 = 10,617,822_{10} \end{aligned}$$

# Converting from decimal to hex

- Use the same idea as for binary
- Example: convert  $1237_{10}$ 
  - $1237 = 77 * 16 + 5$
  - $77 = 4 * 16 + 13$
  - $4 = 0 * 16 + 4$
  - Result:  $4D5_{16}$

# Hexadecimal addition

$$\begin{array}{r} \text{c} \\ \text{A 2 3 F} \\ + \text{3 D 1 3} \\ = \text{D F 5 2} \end{array} \qquad \begin{array}{r} 41535_{10} \\ + 15635_{10} \\ = 57170_{10} \end{array}$$

$$\begin{array}{r} \text{c} \quad \text{c} \quad \text{c} \quad \text{c} \\ \text{D 1 F F} \\ + \text{A 4 D F} \\ = \text{1 7 6 D E} \end{array} \qquad \begin{array}{r} 53759_{10} \\ + 42207_{10} \\ = 95965_{10} \end{array}$$



# Why is hexadecimal useful?

- We need to think in binary because computers operate on binary quantities
- But binary is cumbersome
- However, **hexadecimal makes it possible to represent binary quantities in a compact form**
- Conversions back and forth from binary to hex are straightforward
  - Just convert hex digits into 4-bit numbers
  - Just convert 4-bit binary numbers into hex digits

# Converting from hex to binary

- Consider  $A43FE2_{16}$
- We convert each hex digit into a 4-bit binary number:
  - $A_{16}: 1010_2$
  - $4_{16}: 0100_2$
  - $3_{16}: 0011_2$
  - $F_{16}: 1111_2$
  - $E_{16}: 1110_2$
  - $2_{16}: 0010_2$
- We “glue” them all together:
  - $A43FE2_{16} = 101001000011111111100010_2$
- Note that:
  - **You must have the leading 0's for the 4-bit numbers**, which is what a computer would store anyway
  - It all works because  $F_{16} = 15_{10}$ , and a 4-bit number has maximum value of  $2^4 - 1 = 15_{10}$

# Converting from binary to hex

- Let's convert  $1001010101111_2$  into hex
- We split it in 4-bit numbers, which we convert separately
- First **we add leading 0's** to have a number of bits that's a multiple of 4:

0001 0010 1010 1111

- Then we convert
  - $0001_2 : 1_{16}$
  - $0010_2 : 2_{16}$
  - $1010_2 : A_{16}$
  - $1111_2 : F_{16}$
- And the result:  $1001010101111_2 = 12AF_{16}$



# Integer representation

- A computer needs to store integers in memory/registers
- Stored using different numbers of bytes (1 byte = 8 bits):
  - 1-byte: “byte”
  - 2-byte: “half word” (or “word”)
  - 4-byte: “word” (or “double word”)
  - 8-byte: “double word” (or “paragraph”, or “quadword”)
  - Different computers have used different word sizes, so it’s always a bit confusing to just talk about a “word” without any context
- Regardless of the number of bytes, integers are stored in binary
- Integers come in two flavors:
  - **Unsigned**: values from 0 to  $2^b-1$
  - **Signed**: negatives values, with about the same number of negative values as the number of positive values
- You can actually declare variables as signed or unsigned in some high-level programming languages, like C

# Sign-Magnitude

- **Storing unsigned integers is easy:** just store the bits of the integer's binary representation
- **Storing signed integer raises a question:** how to store the sign?
- One approach is called **sign-magnitude**: reserve the leftmost bit to represent the sign

00100101 denotes  $+0100101_2$

10100101 denotes  $-0100101_2$

- It's very easy to negate a number: just flip the leftmost bit
- Unfortunately, sign-magnitude complicates the logic of the CPU (i.e., ICS331-type stuff)
  - There are two representations for zero: 10000000 and 00000000
  - Some operations are thus more complicated to implement in hardware

# One's complement

- Another idea to store a negative number is to take the complement (i.e., flip all bits) of its positive counterpart
- Example: I want to store integer -87
  - $87_{10} = 01010111_2$
  - $-87_{10} = 10101000$
- Simple, but still two representations for zero: 00000000 and 11111111
- It turns out that computer logic to deal with 1's complement arithmetic is complicated
- Note: it's easy to compute the 1's complement of a number represented in hexadecimal
  - let's consider:  $57_{16}$
  - Subtract each hex digit from F:  $F-5=A$ ,  $F-7=8$
  - 1's complement of  $57_{16}$  is  $A8_{16}$

# Two's complement

- While sign-magnitude and 1's complement were used in older computers, **nowadays all computers use 2's complement**
- Computing the 2's complement is in **two steps**:
  - Compute the 1's complement of the positive number
  - Add 1 to the result
  - This gives the representation of the negative number
- Example: Let's represent  $-87_{10}$ 
  - $87_{10} = 01010111_2$  or  $57_{16}$
  - 1's complement:  $10101000$  or  $A8$
  - Add one:  $10101001$  or  $A9$
- Let's invert again
  - We start with  $A9$
  - Invert:  $56$
  - Add one:  $57$ , which represents  $87_{10}$

# Two's complement

- Note that when adding 1 in the second step a carry may be generated but is ignored!
  - Difference between **arithmetic** and **computer arithmetic**
  - When adding two X-bit quantities in a computer one always obtain another X-bit quantity (X=8, 16, 32, ...)
- Example: Computing 2's complement of 00000000
  - Take the invert: 11111111
  - Add one: 00000000 with a carry generated!
    - Should be a 9-bit quantity: 100000000
- Therefore 0 has only one representation: a signed byte can store values from -128 to +127 (128 < 0 values, and 128 >= 0 values)
- It turns out that 2's complement makes for very simple arithmetic logic when building ALUs
- **From now on we always assumed 2's complement representation**
- **Important:** The leftmost bit still indicates the sign of the number (0: positive, 1: negative)
  - In hex, if the left-most "digit" is 8, 9, A, B, C, D, E, or F, then the number is negative, otherwise it is positive

# Ranges of Numbers

- For 1-byte values
  - Unsigned
    - Smallest value: 00 (0<sub>10</sub>)
    - Largest value: FF (255<sub>10</sub>)
  - Signed
    - Smallest value: 80 (-128<sub>10</sub>)
    - Largest value: 7F (+127<sub>10</sub>)
- For 2-byte values
  - Unsigned
    - Smallest value: 0000 (0<sub>10</sub>)
    - Largest value: FFFF (65,535<sub>10</sub>)
  - Signed
    - Smallest value: 8000 (-32,768<sub>10</sub>)
    - Largest value: 7FFF (+32,767<sub>10</sub>)
- etc.

# The Task of the (Assembly) Programmer

- The computer simply stores data as bits
- The computer internally has no idea what the data means
  - It doesn't know whether numbers are signed or unsigned
- We, as programmers have precise interpretations of what bits mean
  - "I store a 4-byte signed integer", "I store a 1-byte integer which is an ASCII code"
- When using a high-level language like C, we say what data means
  - "I declare x as an int and y as an unsigned char"
- **But** when writing assembly code, we don't have a notion of "data types"
- The ISA provides many instructions that operate on all types of data
- It's our role to use the instructions that correspond to the data
  - e.g., if you used the "signed multiplication" instruction on unsigned numbers, you'll just get a wrong results but no warning/error
- This is one of the difficulties of assembly programming
- And 2's complement appears "magic"...

# The Magic of 2's Complement

- Say I have two 1-byte values, A3 and 17, and I add them together:  
 $A3_{16} + 17_{16} = BA_{16}$
- If my interpretation of the numbers is **unsigned**:
  - $A3_{16} = 163_{10}$
  - $17_{16} = 23_{10}$
  - $BA_{16} = 186_{10}$
  - and indeed,  $163_{10} + 23_{10} = 186_{10}$
- If my interpretation of the numbers is **signed**:
  - $A3_{16} = -93_{10}$
  - $17_{16} = 23_{10}$
  - $BA_{16} = -70_{10}$
  - and indeed,  $-93_{10} + 23_{10} = -70_{10}$
- So, as long as I stick to my interpretation, the **binary addition** does the right thing assuming 2's complement notation!!!
  - Same thing for the subtraction





# Conclusion

- We'll come back to numbers and arithmetic when we use arithmetic assembly instructions